An Algorithm for Stage Semantics

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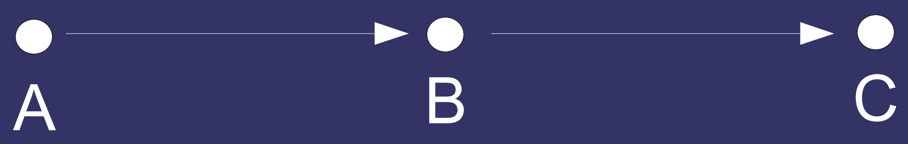

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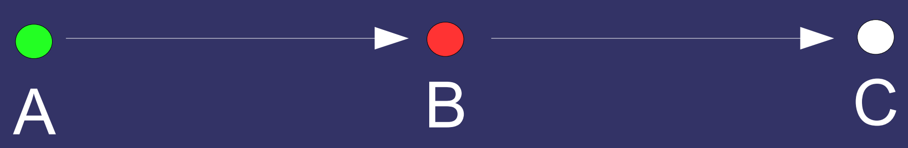

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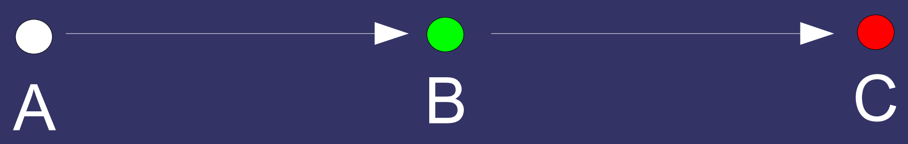

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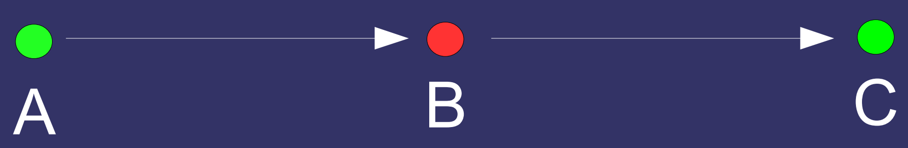

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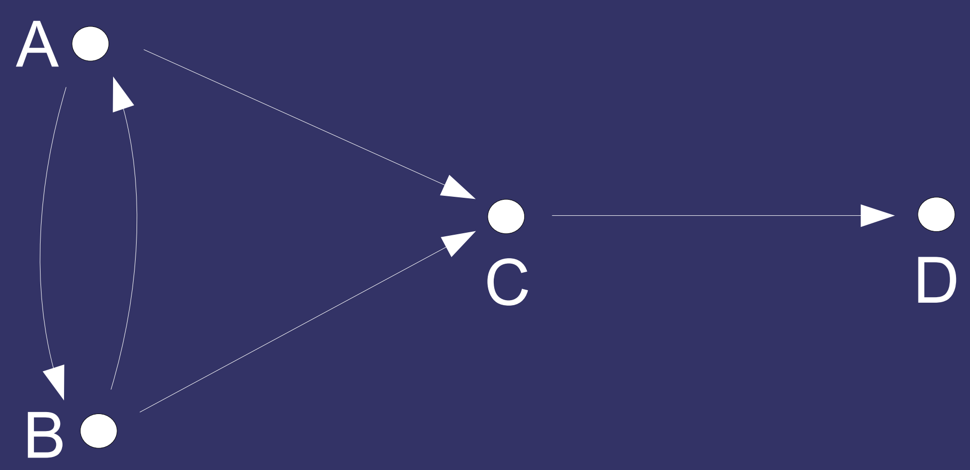

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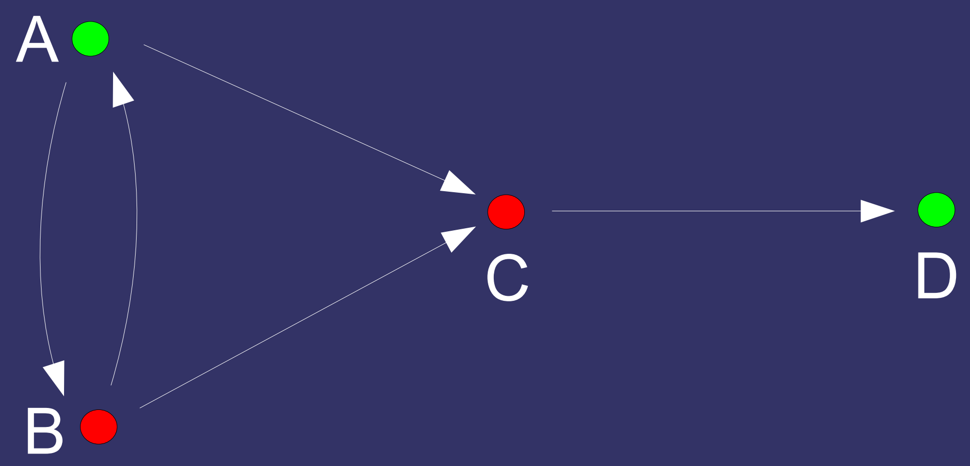

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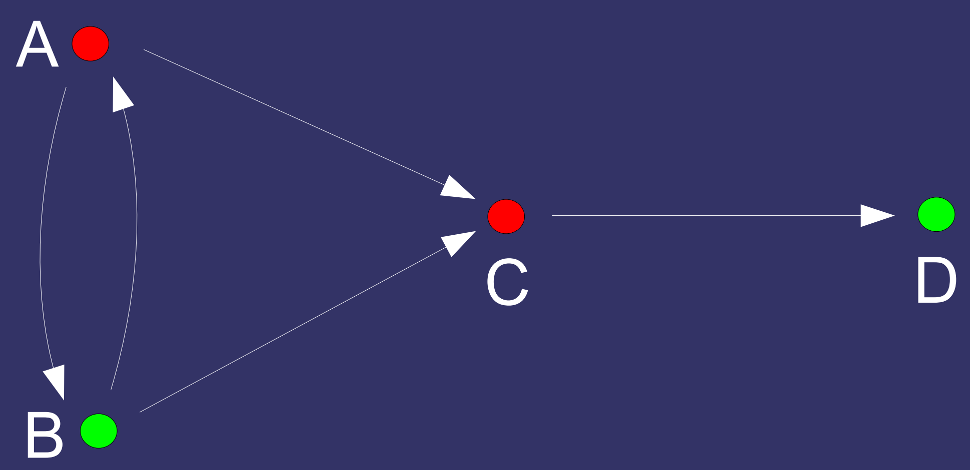

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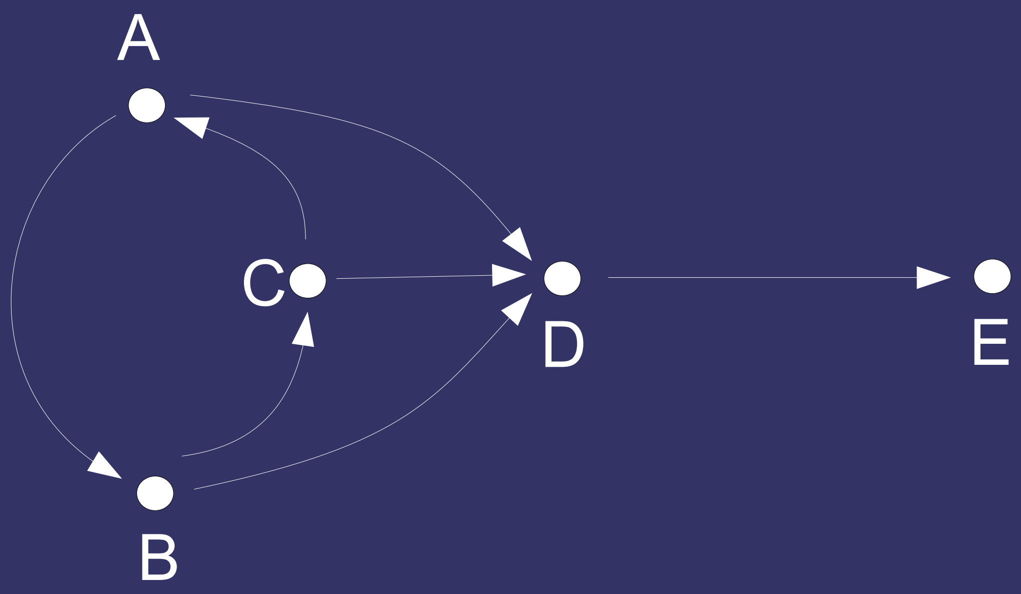

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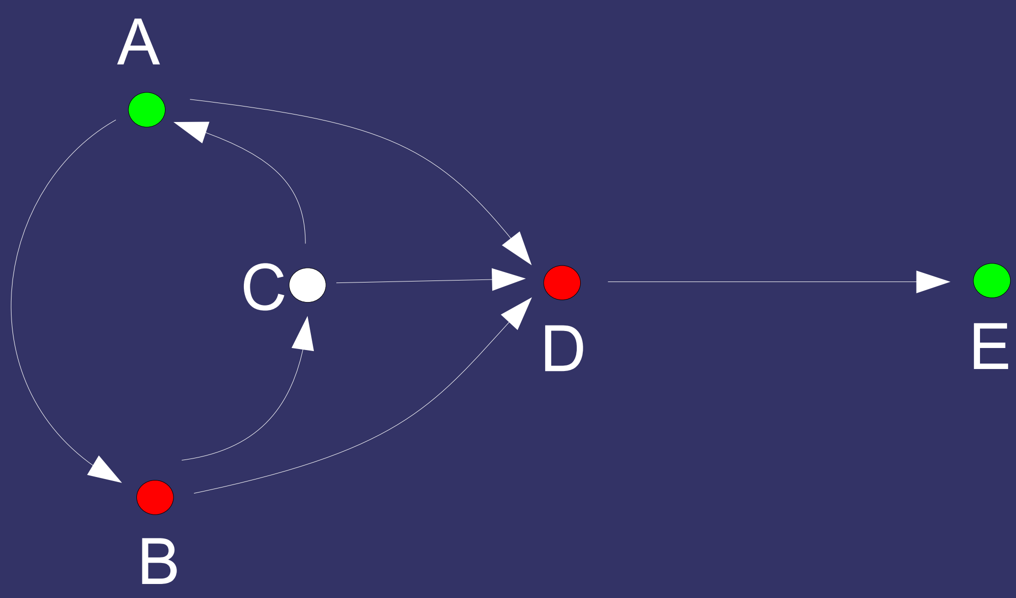

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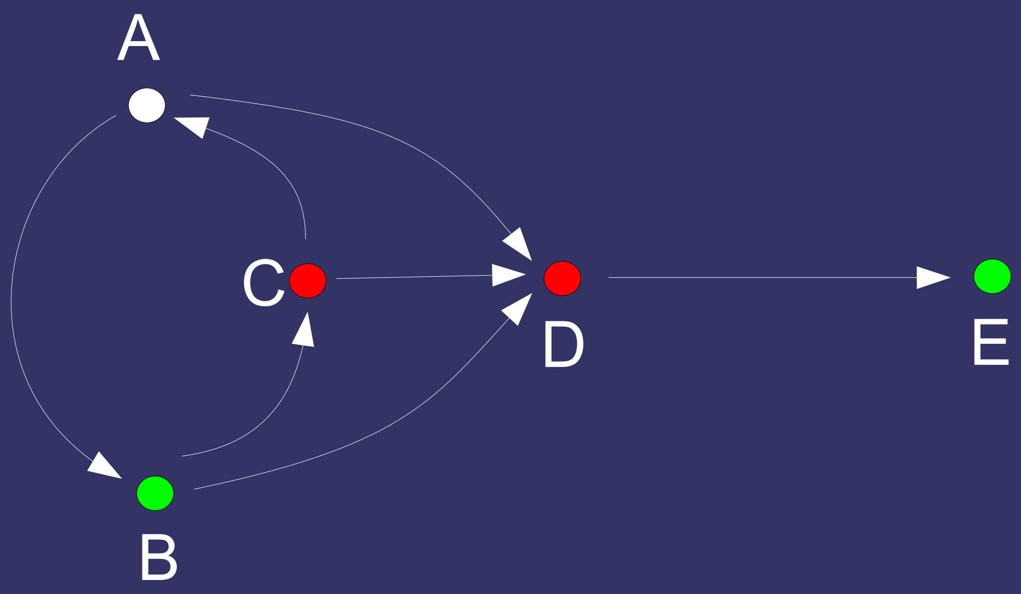

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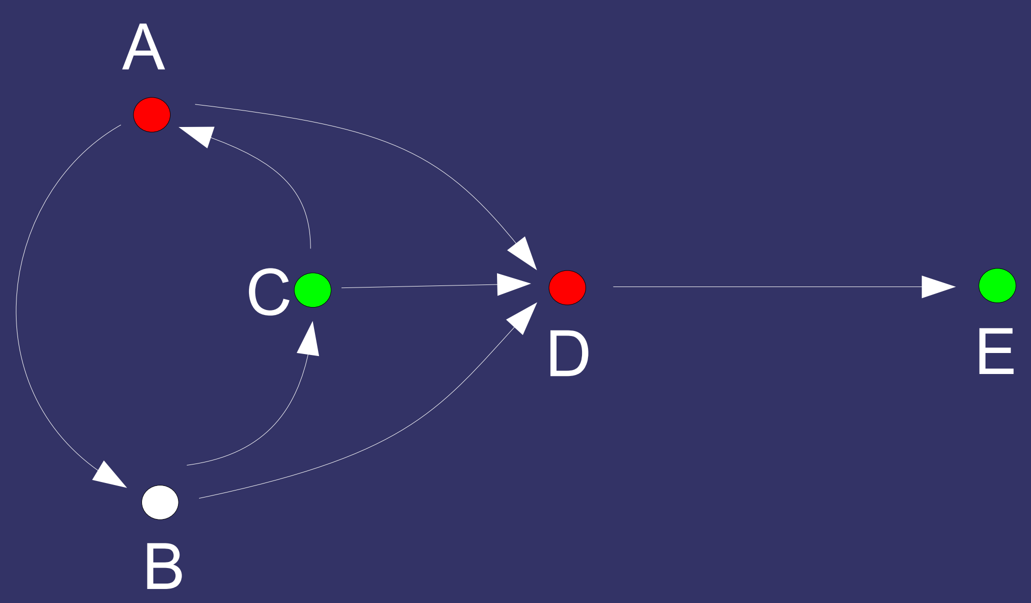

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If there exists a stable extension then every stage extension is also a stable extension

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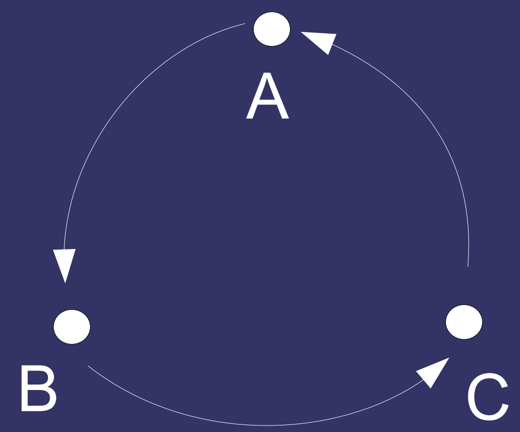

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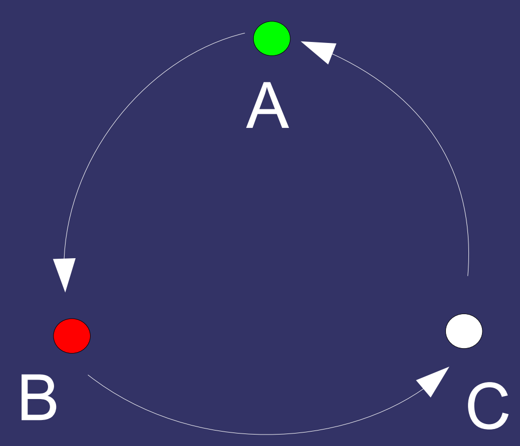

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The “No Models” problem

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When using a semantics that sometimes yields no models (extensions) there are two approaches:

(1) apply the original semantics to maximal subsets of the original problem description

(2) invent a new semantics that can deal with a wider range of problem descriptions

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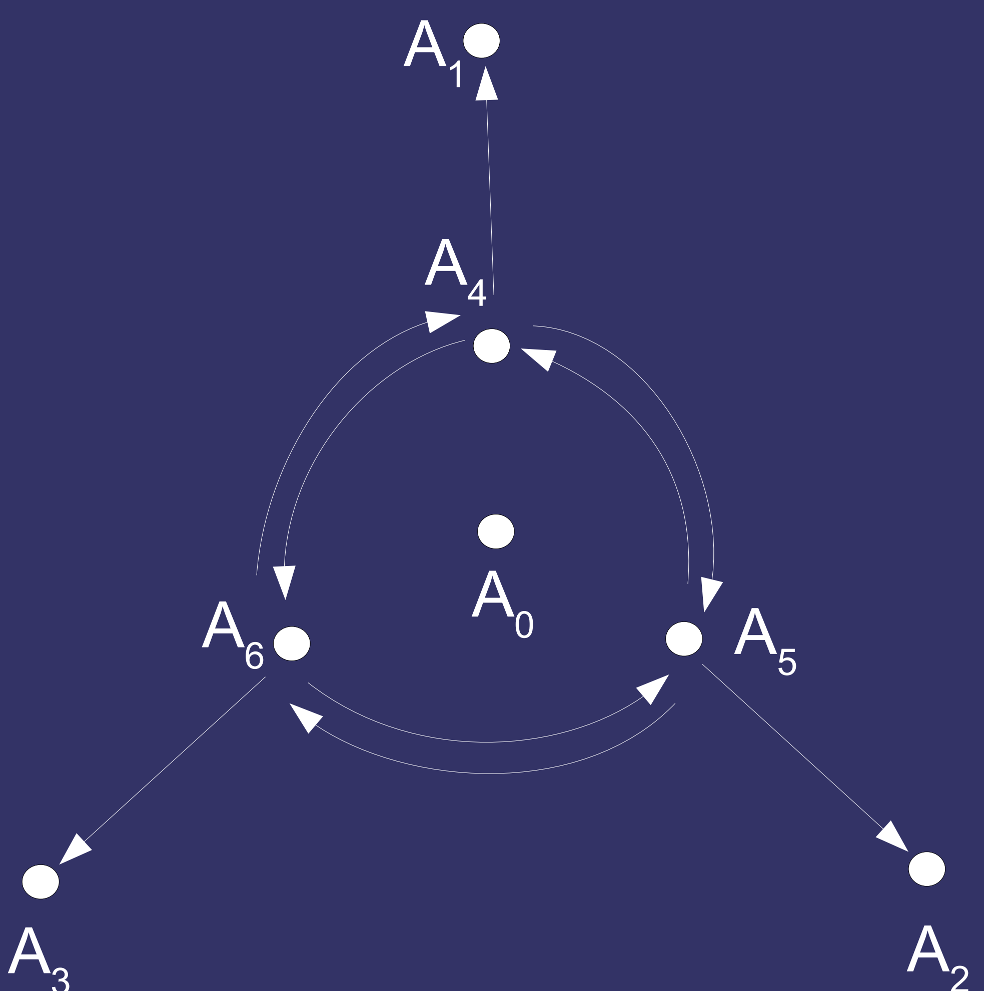

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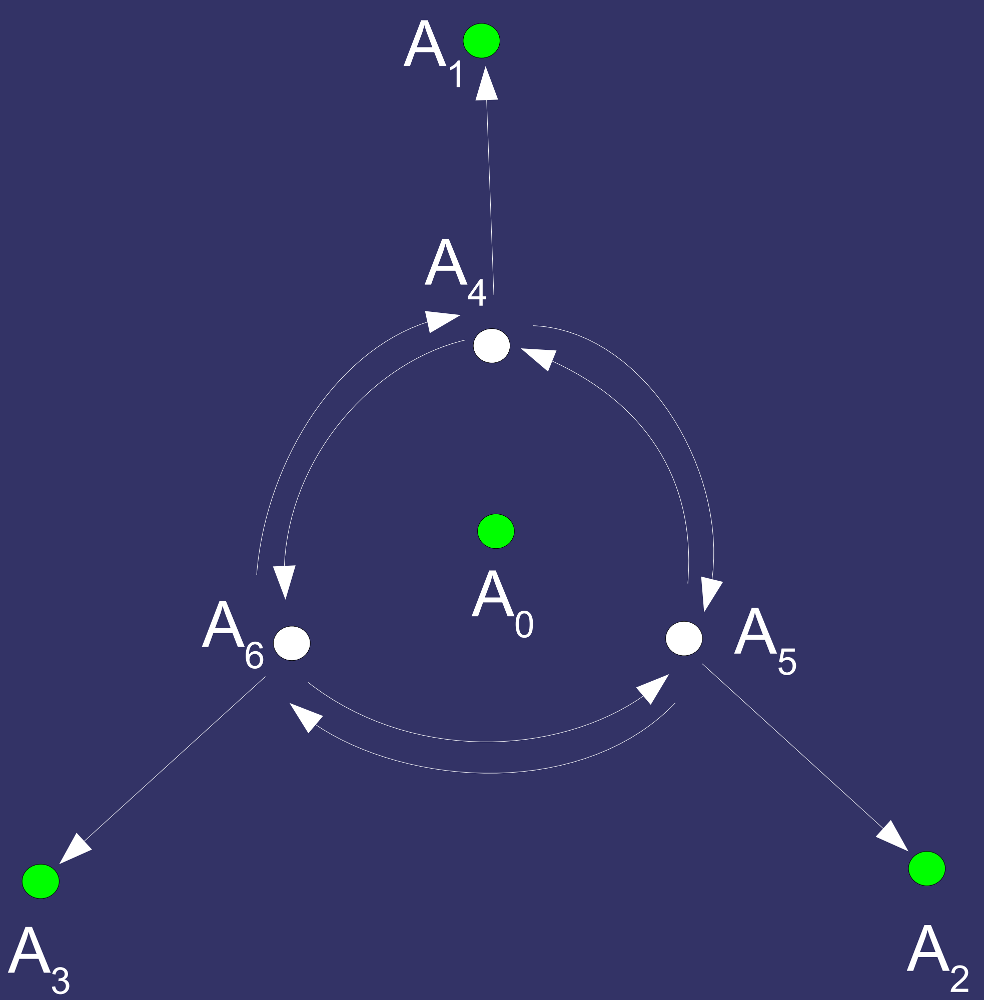

Take Home Message (1/2)

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admissibility ⇒ consistency
(+ constructing the AF the right way) (Caminada & Amgoud, AIJ 2007)

conflict-freeness ⇒ ¬consistency
(+ constructing the AF the same way)

stage semantics ⇒ ¬consistency
(+ constructing the AF the same way)

other non-adm sem. (CF2) ⇒ ?consistency
(+ constructing the AF in any reasonable way)

Take Home Message (2/2)

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Are there any reasonable instantiated argumentation systems that our community's abstract work (CF2, DefLog, ADF, …) provides an abstraction of?

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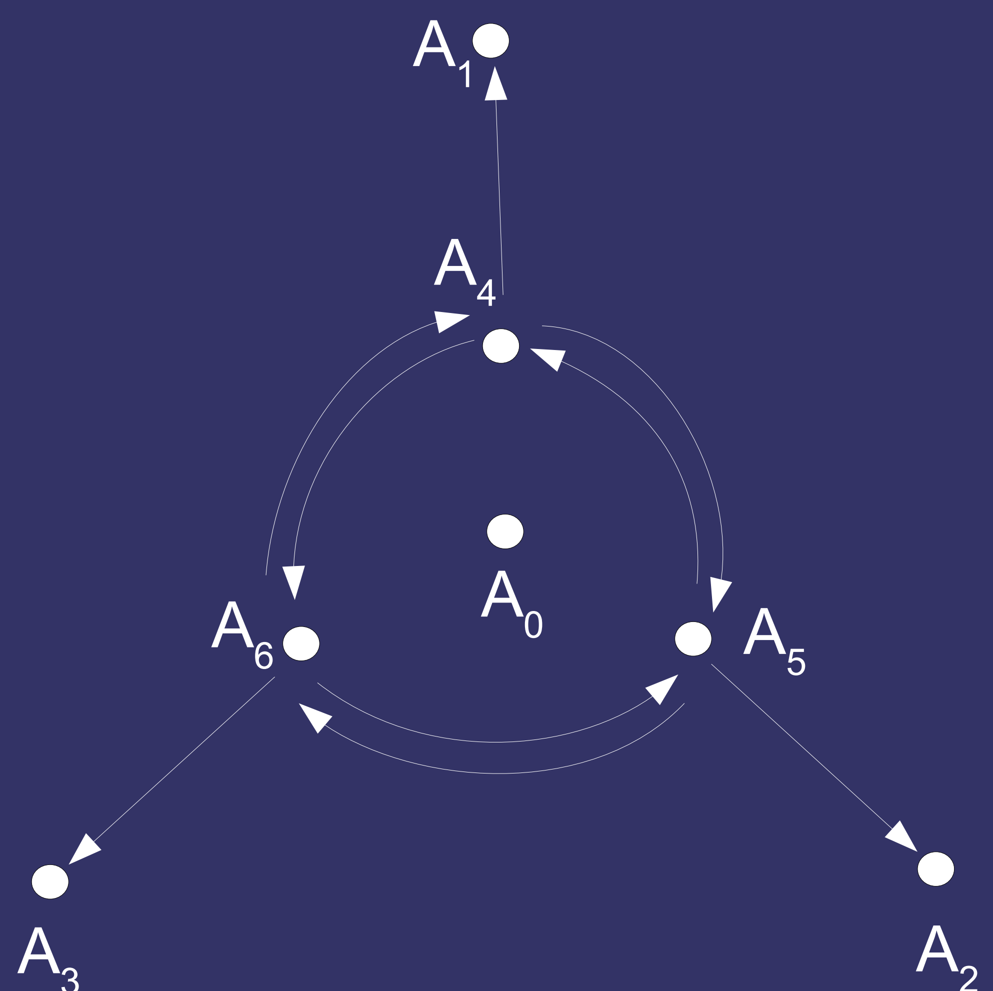

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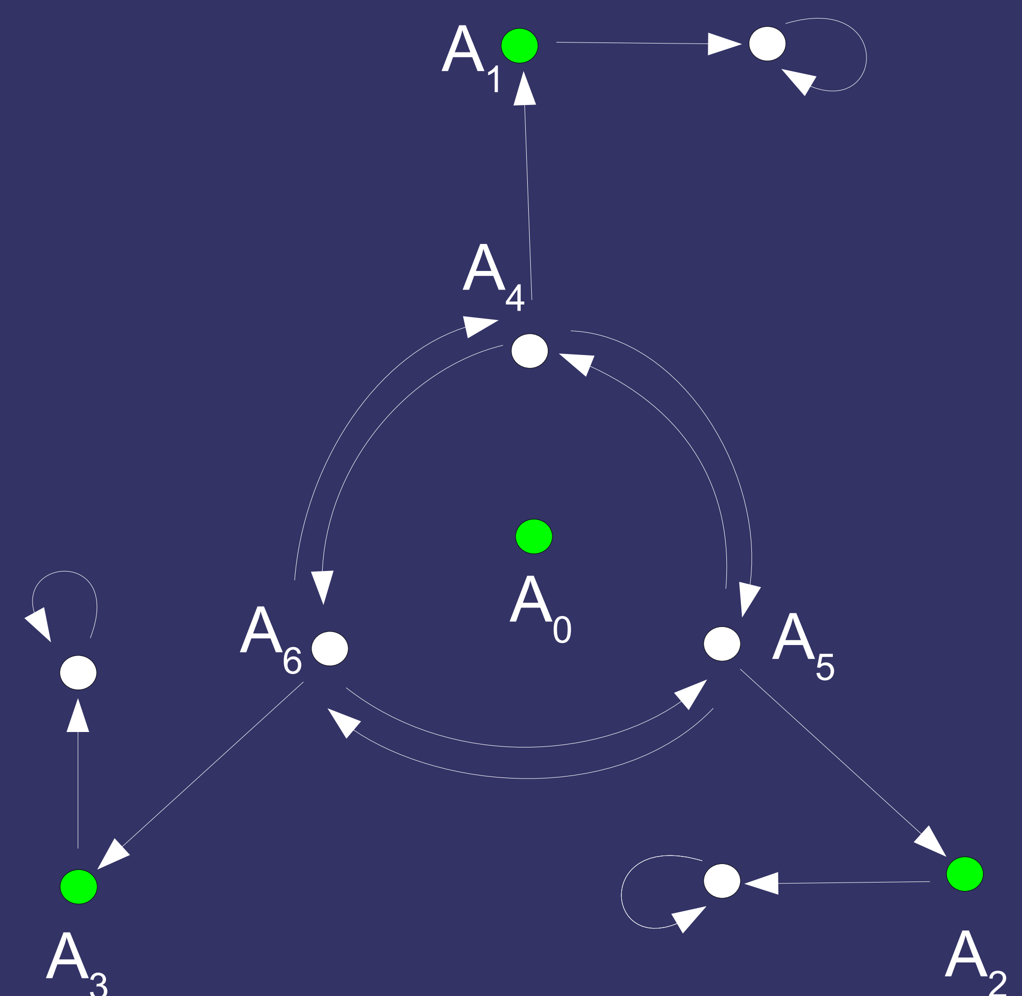

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